

High Availability Solutions

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Load Balancing

Introduction

Network load balancing is the ability to balance traffic across two or more links without using dynamic routing protocols.

There are two type of balancing methods:

- per-packet - each packet of a single stream can be forwarded over different links. This method will work reliably especially on TCP and secure connections only when you are able to control both balancing endpoints.
- per-connection - all packets of the same connection (stream) is always sent over one link. This method is mandatory in setups where only one end of the balancing is under our control, for example, home router with multiple WAN connections.

Method		Per-connection	Per-packet
Firewall Mangle	Nth	Yes	Yes
	PCC (Per Connection Classifier)	Yes	No
	Other matchers	Yes	Yes
ECMP (Equal Cost Multi-Path)		Yes	No
Bonding		No	Yes
OSPF		Yes	No
BGP		Yes	No

Simple Failover Example

Simplest failover setup would be to use multiple gateways when one gateway is active and another one takes over when the first one fails.

To make this work, configure larger **distance** value for the secondary one, and **check-gateway** for the first one:

```
/ip route add gateway=192.168.1.1 distance=1 check-gateway=ping /ip route add gateway=192.168.2.1 distance=2
```

The *check-gateway* will make sure the gateway is up only when actual traffic can reach the gateway. When the ping fails the first gateway will become inactive and the second one will take over, and when the first gateway recovers it will become active and make the second gateway to work again as a backup.

Per connection classifier

PCC matcher will allow you to divide traffic into equal streams with the ability to keep packets with specific set of options in one particular stream (you can specify this set of options from src-address, src-port, dst-address, dst-port)

Theory

PCC takes selected fields from IP header, and with the help of a hashing algorithm converts selected fields into 32-bit value. This value then is divided by a specified *Denominator* and the remainder then is compared to a specified *Remainder*, if equal then the packet will be captured. You can choose from src-address, dst-address, src-port, dst-port from the header to use in this operation.

PCC is not a valid method, in case of Hotspot(captive portal) - due to the fact that Hotspot uses web-proxy and it uses only the default routing table, at the moment.

```
per-connection-classifier=  
PerConnectionClassifier ::= [!]ValuesToHash:Denominator/Remainder  
Remainder ::= 0..4294967295 (integer number)  
Denominator ::= 1..4294967295 (integer number)  
ValuesToHash ::= both-addresses|both-ports|dst-address-and-port|  
src-address|src-port|both-addresses-and-ports|dst-address|dst-port|src-address-and-port
```

Example

This configuration will divide all connections into 3 groups based on the source address and port

```
/ip firewall mangle add chain=prerouting action=mark-connection \  
new-connection-mark=1st_conn per-connection-classifier=src-address-and-port:3/0  
/ip firewall mangle add chain=prerouting action=mark-connection \  
new-connection-mark=2nd_conn per-connection-classifier=src-address-and-port:3/1
```

```
/ip firewall mangle add chain=prerouting action=mark-connection \
new-connection-mark=3rd_conn per-connection-classifier=src-address-and-port:3/2
```

How PCC works

This article aims to explain in simple terms how PCC works. The definition from the official manual wiki page reads: "PCC takes selected fields from IP header, and with the help of a hashing algorithm converts selected fields into 32-bit value. This value then is divided by a specified Denominator and the remainder then is compared to a specified Remainder, if equal then packet will be captured. You can choose from src-address, dst-address, src-port, dst-port from the header to use in this operation.", with the full number of fields available being: "both-addresses|both-ports|dst-address-and-port|src-address|src-port|both-addresses-and-ports|dst-address|dst-port|src-address-and-port". If you understand that definition, there'll be nothing interesting in this article for you.

First, here are the terms necessary to understand the definition.

IP packets have a header that contains several fields, two of those fields are the IP address of the source of the packet and the IP address of the destination of the packet. TCP and UDP packets also have headers that contain the source port and the destination port.

Denominators and remainders are parts of modulus operations. A modulus operation produces the integer left over when you divide two numbers and only accept the whole number portion of the result. It is represented by a % sign. Here are some examples: $3 \% 3 = 0$, because 3 divides cleanly by 3. $4 \% 3 = 1$, because the next smallest number to 4 that cleanly divides by 3 is 3, and $4 - 3 = 1$. $5 \% 3 = 2$, because the next smallest number to 5 that divides cleanly by 3 is 3, and $5 - 3 = 2$. $6 \% 3 = 0$, because 6 divides cleanly by 3.

A hash is a function that is fed input, and produces output. Hashes have many interesting properties, but the only important one for the purpose of this article is that hash functions are deterministic. That means that when you feed a hash function an input that reads 'hello' and it produces the output '1', you can rely on the fact that if you feed it 'hello' a second time it will produce the output '1' again. When you feed a hash function the same input, it will always produce the same output. What exact hashing algorithm is used by PCC is not important, so for this discussion let's assume that when you feed it IP addresses and ports, it just adds up the octets of the IP addresses as decimal numbers as well as the ports, and then takes the last digit and produces it as the output. Here an example:

The hash function is fed 1.1.1.1 as the source IP address, 10000 as the source TCP port, 2.2.2.2 as the destination IP address and 80 as the destination TCP port. The output will be $1+1+1+1+10000+2+2+2+2+80 = 10092$, the last digit of that is 2, so the hash output is 2. It will produce 2 every time it is fed that combination of IP addresses and ports.

At this point it's important to note that even though PCC is most often used for spreading load across circuits, PCC itself has absolutely nothing to do with routing, routing marks or spreading load. PCC is simply a way to match packets, and not directly related to the action of then marking those matched packets even if that is its main purpose.

Here are three lines often used for PCC, with their explanation:

```
/ip firewall mangle add chain=prerouting action=mark-connection \
new-connection-mark=1st_conn per-connection-classifier=src-address-and-port:3/0
/ip firewall mangle add chain=prerouting action=mark-connection \
new-connection-mark=2nd_conn per-connection-classifier=src-address-and-port:3/1
/ip firewall mangle add chain=prerouting action=mark-connection \
new-connection-mark=3rd_conn per-connection-classifier=src-address-and-port:3/2
```

Here are what the different field options mean for the purpose of packet matching, these are the fields that will be fed into the hashing algorithm (and, for the purpose of spreading load across links, decide what link a packet will be put on). Remember that a hash function will always produce the same input when it's fed the same output:

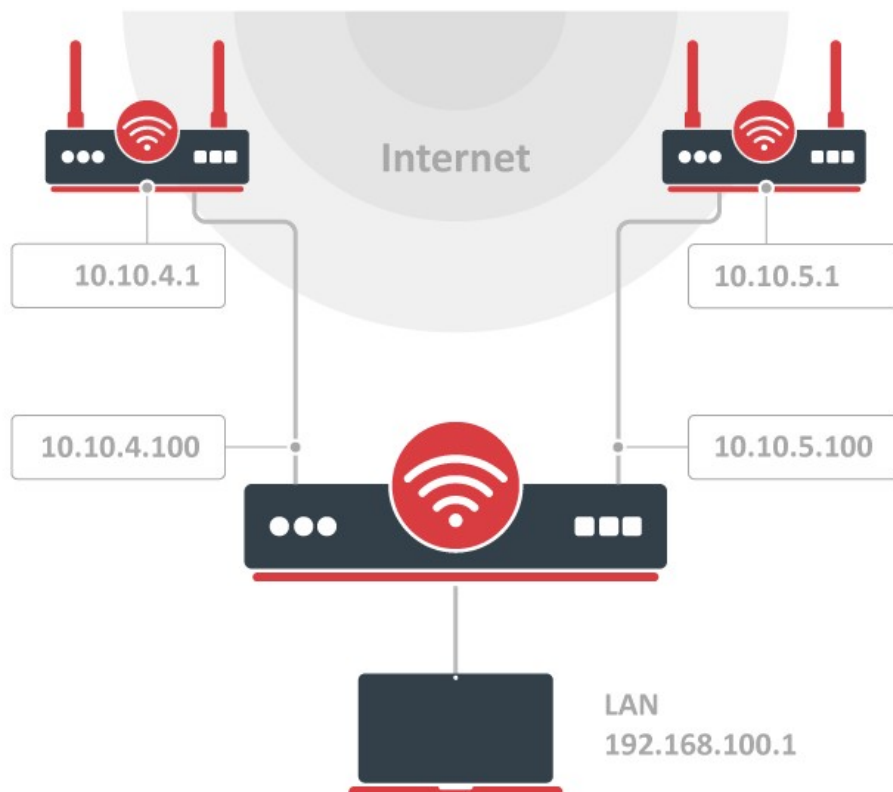
- **src-address:** The source address of a client will always be the same, so all traffic from a particular client will always match the same PCC matcher, and will always be put on the same link.
- **dst-address:** The destination address of a specific server will always be the same, so all traffic to that server (say, the Mikrotik Wiki) will always match the same PCC matcher, and will always be put on the same link.
- **both-addresses:** The source and destination IP pair between the same client and server will always be the same, so all traffic between one specific client and a specific server (say, your laptop and the Mikrotik Wiki) will always match the same PCC matcher, and will always be put on the same link.
- **src-port:** Source ports of clients are usually randomly chosen when the connection is created, so across many connections different source ports will be fed into the hash function, and different PCC matchers will match and traffic will go across different links. However, some client protocols always choose the same source port, and servers behind your router will mostly likely always use the same service port to send traffic back to their clients. A web server behind your router would send most traffic from its HTTP (80) and HTTPS (443) ports, and that traffic would always match the same PCC matcher and would be put on the same link.
- **dst-port:** Destination ports of clients are usually well defined service ports, all HTTP (80) traffic between your clients and servers on the Internet would always match the same PCC matcher, and would be put on the same link. However, the same clients doing HTTPS (443) traffic could match a different PCC matcher, and would go across a different link.
- **both-ports:** Since the client port is (usually) randomly chosen, the combination of the two ports is (usually) random and will spread load across links.
- **src-address-and-port:** Same caveat as src-port.
- **dst-address-and-port:** Same caveat as dst-port.

- both-addresses-and-ports: This is the most random way to spread traffic across links, since it has the most number of variables.

It's important to note that even though the hash function discussed in this article is greatly simplified and not what is used in real life, it nicely demonstrates another property of hash functions: two completely different inputs can produce the same output. In our example, $3 \% 3 = 0$, and $6 \% 3 = 0$; we get back 0 when we feed it a 3 as well as when we feed it a 6. The same is true for the actual function used for PCC, even though I don't know what it is we do know from the definition that it produces a 32 bit value as output. IP addresses are 32 bit, and ports are 16 bit, so assuming that we're using both-addresses-and-ports, we'd be feeding it $32+32+16+16 = 96$ bits of input and would only receive 32 bits back, so it must be producing the same output for different inputs. This means that two completely unrelated connections could match the same PCC matcher, and would be put on the same line. PCC works better the more connections you put across it so that the hash function has more chances to produce different outputs.

Configuration Example

Let's assume this configuration:



IP Addresses

```
/ip address add address=10.10.4.100/24 interface=ether_ISP1 network=10.10.4.0
add address=10.10.5.100/24 interface=ether_ISP2 network=10.10.5.0
add address=192.168.100.1/24 interface=ether_LAN network=192.168.100.0
```

The router has two upstream (ISP) interfaces with the addresses of 10.10.4.100/24 and 10.10.5.100/24. The LAN interface has IP address of 192.168.0.1/24.

We are adding two new Routing tables, which will be used later:

```
/routing table
add disabled=no fib name=ISP1_table
add disabled=no fib name=ISP2_table
```

Policy routing

```
/ip firewall mangle
add action=accept chain=prerouting dst-address=10.10.4.0/24 in-interface=ether_LAN
add action=accept chain=prerouting dst-address=10.10.5.0/24 in-interface=ether_LAN
```

With policy routing it is possible to force all traffic to the specific gateway, even if traffic is destined to the host (other than gateway) from the connected networks. This way routing loop will be generated and communications with those hosts will be impossible. To avoid this situation we need to allow usage of default routing table for traffic to connected networks.

```
add action=mark-connection chain=input connection-state=new in-interface=ether_ISP1 new-connection-mark=ISP1
add action=mark-connection chain=input connection-state=new in-interface=ether_ISP2 new-connection-mark=ISP2

add action=mark-connection chain=output connection-mark=no-mark connection-state=new new-connection-mark=ISP1 passthrough=yes per-connection-classifier=both-addresses:2/0
add action=mark-connection chain=output connection-mark=no-mark connection-state=new new-connection-mark=ISP2 per-connection-classifier=both-addresses:2/1
```

First it is necessary to manage connection initiated from outside - replies must leave via same interface (from same Public IP) request came. We will mark all new incoming connections, to remember what was the interface.

```
add action=mark-connection chain=prerouting connection-mark=no-mark connection-state=new dst-address-type=!local in-interface=ether_LAN new-connection-mark=ISP1 per-connection-classifier=both-addresses:2/0
add action=mark-connection chain=prerouting connection-mark=no-mark connection-state=new dst-address-type=!local in-interface=ether_LAN new-connection-mark=ISP2 per-connection-classifier=both-addresses:2/1
```

Action mark-routing can be used only in mangle chain output and prerouting, but mangle chain prerouting is capturing all traffic that is going to the router itself. To avoid this we will use `dst-address-type=!local`. And with the help of the new PCC we will divide traffic into two groups based on source and destination addressees.

```
add action=mark-routing chain=output connection-mark=ISP1 new-routing-mark=ISP1_table
add action=mark-routing chain=prerouting connection-mark=ISP1 in-interface=ether_LAN new-routing-mark=ISP1_table
```

```
add action=mark-routing chain=output connection-mark=ISP2 new-routing-mark=ISP2_table
add action=mark-routing chain=prerouting connection-mark=ISP2 in-interface=ether_LAN new-routing-mark=ISP2_table
```

Then we need to mark all packets from those connections with a proper mark. As policy routing is required only for traffic going to the Internet, do not forget to specify in-interface option.

```
/ip route
add check-gateway=ping disabled=no dst-address=0.0.0.0/0 gateway=10.10.4.1 routing-table=ISP1_table
suppress-hw-offload=no
add check-gateway=ping disabled=no dst-address=0.0.0.0/0 gateway=10.10.5.1 routing-table=ISP2_table
suppress-hw-offload=no
```

Create a route for each routing-mark

```
add distance=1 dst-address=0.0.0.0/0 gateway=10.10.4.1
add distance=2 dst-address=0.0.0.0/0 gateway=10.10.5.1
```

To enable failover, it is necessary to have routes that will jump in as soon as others will become inactive on gateway failure. (and that will happen only if check-gateway option is active)

NAT

```
/ip firewall nat
add action=masquerade chain=srcnat out-interface=ether_ISP1
add action=masquerade chain=srcnat out-interface=ether_ISP2
```

As routing decision is already made we just need rules that will fix src-addresses for all outgoing packets. If this packet will leave via ether_ISP1 it will be NATed to 10.10.4.100, if via ether_ISP2 then NATed to 10.10.5.100

Bonding

Summary

Bonding is a technology that allows aggregation of multiple ethernet-like interfaces into a single virtual link, thus getting higher data rates and providing failover.

Interface bonding does not create an interface with a larger link speed. Interface bonding creates a virtual interface that can load balance traffic over multiple interfaces.

CRS3xx, CRS5xx series switches, CCR2116, CCR2216 routers and 88E6393X, 88E6191X, 88E6190 switch chips support bridge hardware offloading with bonding interfaces. Only `802.3ad` and `balance-xor` bonding modes are hardware offloaded, other bonding modes will use the CPU's resources. The built-in switch chip will always use Layer2+Layer3+Layer4 for a transmit hash policy, changing the transmit hash policy manually will have no effect.

Quick Setup Guide

Let us assume that we have two Ethernet interfaces on each router (Router1 and Router2) and want to get the maximum data rate between these two routers. To make this possible, follow these steps:

1. Make sure that you do not have IP addresses on interfaces that will be enslaved for bonding interface.
2. Add bonding interface and IP address on the Router1:

```
/interface bonding add slaves=ether1,ether2 name=bond1
/ip address add address=172.16.0.1/24 interface=bond1
```

3. Do the same thing on the Router2:

```
/interface bonding add slaves=ether1,ether2 name=bond1
/ip address add address=172.16.0.2/24 interface=bond1
```

4. Test the link from Router1:

```
[admin@Router1] > ping 172.16.0.2
```

SEQ	HOST	SIZE	TTL	TIME	STATUS
0	172.16.0.2	56	64	0ms	
1	172.16.0.2	56	64	0ms	
2	172.16.0.2	56	64	0ms	

```
sent=3 received=3 packet-loss=0% min-rtt=0ms avg-rtt=0ms max-rtt=0ms
```

The bonding interface needs a couple of seconds to get connectivity with its peers.

Link monitoring

It is critical that one of the available link monitoring options is enabled. In the above example, if one of the bonded links were to fail, the bonding driver will still continue to send packets over the failed link which will lead to network degradation. Bonding in RouterOS currently supports two schemes for monitoring a link state of slave devices: MII and ARP monitoring. It is not possible to use both methods at the same time due to restrictions in the bonding driver.

ARP Monitoring

ARP monitoring sends ARP queries and uses the response as an indication that the link is operational. The ARP replies are not validated, any received packet by the slave interface will result in the slave interface considered as active. This gives assurance that traffic is actually flowing over the links. If balance-rr and balance-xor modes are set, then the switch should be configured to evenly distribute packets across all links. Otherwise, all replies from the ARP targets will be received on the same link which could cause other links to fail. ARP monitoring is enabled by setting three properties - link-monitoring, arp-ip-targets and arp-interval. The meaning of each option is described later in this article. It is possible to specify multiple ARP targets that can be useful in High Availability setups. If only one target is set, the target itself may go down. Having additional targets increases the reliability of the ARP monitoring.

To enable ARP monitoring on Router1:

```
/interface bonding set [find name=bond1] link-monitoring=arp arp-ip-targets=172.16.0.2
```

and Router2:

```
/interface bonding set [find name=bond1] link-monitoring=arp arp-ip-targets=172.16.0.1
```

We will not change the arp-interval value in our example, RouterOS sets arp-interval to 100ms by default. Unplug one of the cables to test if the link monitoring works correctly, you might notice some ping timeouts until arp monitoring detects link failure.

```
[admin@MikroTik] > ping 172.16.0.2
```

SEQ	HOST	SIZE	TTL	TIME	STATUS
0	172.16.0.2	56	64	0ms	
1	172.16.0.2	56	64	0ms	
2	172.16.0.2	56	64	0ms	
3	172.16.0.2	56	64	0ms	
4	172.16.0.2				timeout
5	172.16.0.2	56	64	0ms	
6	172.16.0.2	56	64	0ms	

```
sent=7 received=6 packet-loss=14% min-rtt=0ms avg-rtt=0ms max-rtt=0ms
```

For ARP monitoring to work properly it is not required to have any IP address on the device, ARP monitoring will work regardless of the IP address that is set on any interface.

When ARP monitoring is used, bonding slaves will send out ARP requests without a VLAN tag, even if an IP address is set on a VLAN interface in the same subnet as the arp-ip-targets

MII monitoring

MII monitoring monitors only the state of the local interface. *MII Type 1* - a device driver determines whether a link is up or down. If the device driver does not support this option then the link will appear as always up. The main disadvantage is that MII monitoring can't tell if the link can actually pass packets or not, even if the link is detected as being up. MII monitoring is configured by setting the variables - link-monitoring and mii-interval.

To enable MII Type1 monitoring on Router1 and Router2:

```
/interface bonding set [find name=bond1] link-monitoring=mii
```

We will leave mii-interval to its default value (100ms). When unplugging one of the cables, the failure will be detected almost instantly compared to ARP link monitoring.

Bonding modes

802.3ad

802.3ad mode is an IEEE standard also called LACP (Link Aggregation Control Protocol). It includes automatic configuration of the aggregates, so minimal configuration of the switch is needed. This standard also mandates that frames will be delivered in order and connections should not see misordering of packets. The standard also mandates that all devices in the aggregate must operate at the same speed and duplex mode.

LACP balances outgoing traffic across the active ports based on hashed protocol header information and accepts incoming traffic from any active port. The hash includes the Ethernet source and destination address and if available, the VLAN tag, and the IPv4/IPv6 source and destination address. How this is calculated depends on transmit-hash-policy parameter. The ARP link monitoring is not recommended, because the ARP replies might arrive only on one slave port due to transmit hash policy on the LACP peer device. This can result in unbalanced transmitted traffic, so MII link monitoring is the recommended option.

The layer-3-and-4 transmit hash mode is not fully compatible with LACP. More details can be found in <https://www.kernel.org/doc/Documentation/networking/bonding.txt>

balance-xor

This mode balances outgoing traffic across the active ports based on the hashed protocol header information and accepts incoming traffic from any active port. The mode is very similar to [LACP](#) except that it is not standardized and works with **layer-3-and-4** hash policy. The mode can work together with static Link Aggregation Group (LAG) interfaces.

balance-rr

If this mode is set, packets are transmitted in sequential order from the first available slave to the last. The balance-rr is the only mode that will send packets across multiple interfaces that belong to the same TCP/IP connection. When utilizing multiple sending and multiple receiving links, packets are often received out of order, which results in segment retransmission, for other protocols such as UDP it is not a problem if a client software can tolerate out-of-order packets. If a switch is used to aggregate links together, then appropriate switch port configuration is required, however many switches do not support balance-rr. [Quick setup guide](#) demonstrates the usage of the balance-rr bonding mode. As you can see, it is quite simple to set up. Balance-rr is also useful for bonding several wireless links, however, it requires equal bandwidth for all bonded links. If the bandwidth of one bonded link drops, then the total bandwidth of bond will be equal to the bandwidth of the slowest bonded link.

active-backup

This mode uses only one active slave to transmit packets. The additional slave only becomes active if the primary slave fails. The MAC address of the bonding interface is presented onto the active port to avoid confusing the switch. Active-backup is the best choice in high availability setups with multiple switches that are interconnected.

The ARP monitoring in this mode will not work correctly if both routers are directly connected. In such setups, MII monitoring must be used or a switch should be put between routers.

broadcast

When ports are configured with broadcast mode, all slave ports transmit the same packets to the destination to provide fault tolerance. This mode does not provide load balancing.

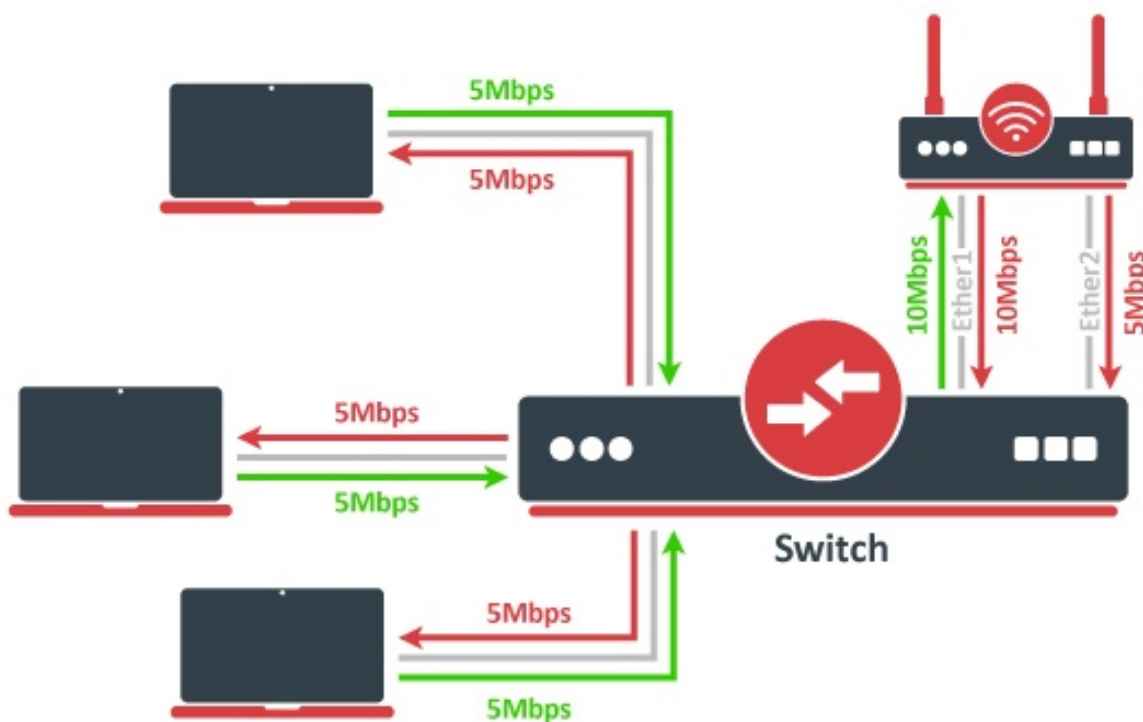
balance-tlb

This mode balances outgoing traffic by peer. Each link can be a different speed and duplex mode and no specific switch configuration is required as for the other modes. The downside of this mode is that only MII link monitoring is supported (ARP link monitoring is ignored when configured) and incoming traffic is not balanced. Incoming traffic will use the link that is configured as "primary".

Configuration example

Let's assume that the router has two links - **ether1** max bandwidth is 10Mbps and **ether2** max bandwidth is 5Mbps. The first link has more bandwidth so we set it as a primary link:

```
/interface bonding add mode=balance-tlb slaves=ether1,ether2 primary=ether1
```

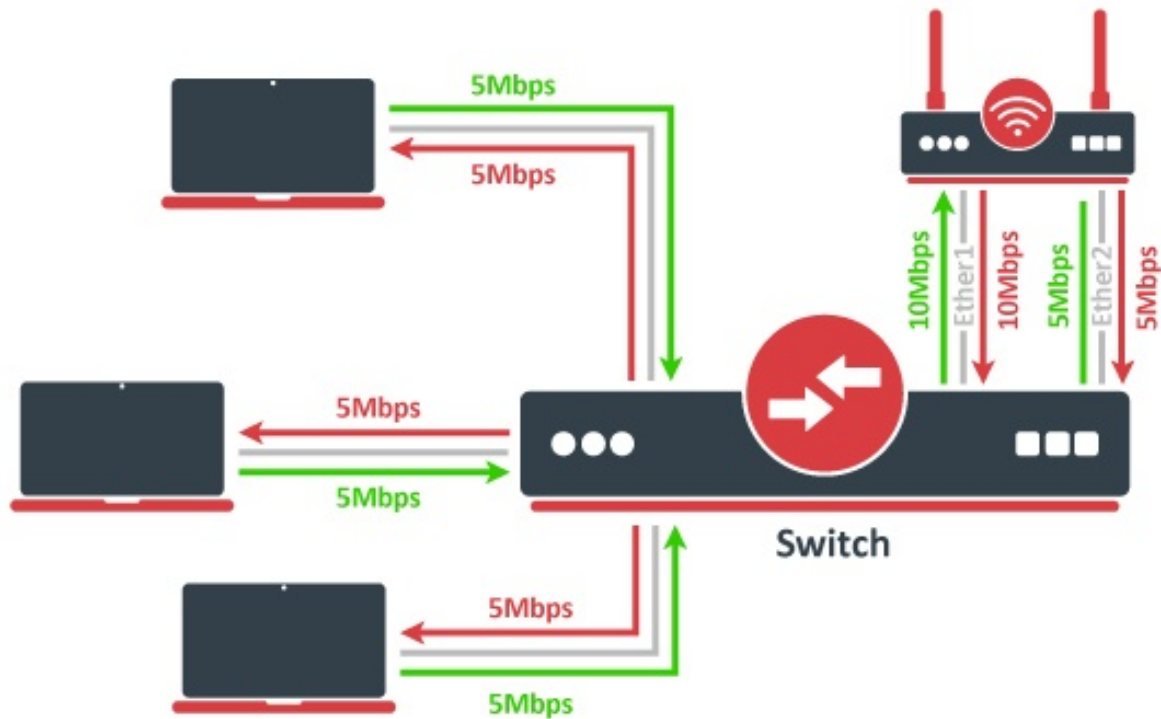


No additional configuration is required for the switch. The image above illustrates how `balance-tlb` mode works. As you can see router can communicate to all the clients connected to the switch with a total bandwidth of both links (15Mbps). But as you already know, `balance-tlb` is not balancing incoming traffic. In our example, clients can communicate to the router with a total bandwidth of primary link which is 10Mbps in our configuration.

balance-alb

The mode is basically the same as `balance-tlb` but incoming IPv4 traffic is also balanced. The receive load balancing is achieved by ARP negotiation. The bonding driver intercepts locally

generated ARP messages on their way out and overwrites the source hardware address with the unique address of one of the slaves in the bond such that different peers use different hardware addresses. Only MII link monitoring is supported (ARP link monitoring is ignored when configured), the additional downside of this mode is that it requires device driver capability to change MAC address. The mode is not compatible with local-proxy-arp setting.



The image above illustrates how balance-alb mode works. Compared to balance-tlb mode, traffic from clients can also use the secondary link to communicate with the router.

Bonding monitoring

Since RouterOS 6.48 version, it is possible to monitor the bonding interface and bonding ports. For the `802.3ad` bonding mode, more detailed monitoring options are available.

```
/interface bonding monitor [find]
    mode: 802.3ad      active-backup
    active-ports: ether4 ether6
                    ether5
    inactive-ports: ether7
    lacp-system-id: CC:2D:E0:11:22:33
    lacp-system-priority: 65535
    lacp-partner-system-id: B8:69:F4:44:55:66
```

Property	Description
mode (802.3ad active-backup balance-alb balance-rr balance-tlb balance-xor broadcast)	Used bonding mode
active-ports (interface)	Shows the active bonding ports
inactive-ports (interface)	Shows the inactive bonding ports (e.g. a disabled or backup interface)
lACP-system-id (MAC address)	Shows the local LACP system ID
lACP-system-priority (integer)	Shows the local LACP priority
lACP-partner-system-id (MAC address)	Shows the partner LACP system ID

To monitor individual bonding ports, use a `monitor-slaves` command.

```
/interface bonding monitor-slaves bond1
```

Flags: A - active, P - partner

```
AP port=ether4 key=17 flags="A-GSCD--" partner-sys-id=D4:CA:6D:12:06:65 partner-sys-priority=65535 partner-key=9 partner-flags="A-GSCD--"
```

```
AP port=ether5 key=17 flags="A-GSCD--" partner-sys-id=D4:CA:6D:12:06:65 partner-sys-priority=65535 partner-key=9 partner-flags="A-GSCD--"
```

Property	Description
port (interface)	Used bonding port
key (integer)	Shows the local LACP aggregation key. The lower 6 bits are automatically assigned based on individual port link speed and duplex. The upper 10 bits can be manually specified using the <code>lACP-user-key</code> setting (available only since RouterOS v7.3).
flags (string)	Shows the local LACP flags: A - activity (link is active, otherwise passive) T - timeout (link is using short 1-second timeout, otherwise using 30-second timeout) G - aggregation (link can be aggregatable) S - synchronization (link is synchronized) C - collecting (link is able to collect incoming frames) D - distributing (link is able to distribute outgoing frames) F - defaulted (link is using defaulted partner information, indicated that no LACPDU has been received from the partner) E - expired (link has expired state)
partner-sys-id (MAC address)	Shows the partner LACP system ID
partner-sys-priority (integer)	Shows the partner LACP priority
partner-key (integer)	Shows the partner LACP aggregation key
partner-flags (string)	Shows the partner LACP flags

Property Description

This section describes the available bonding settings.

Property	Description
arp (<i>disabled enabled proxy-arp reply-only</i> ; Default: enabled)	Address Resolution Protocol for the interface. <ul style="list-style-type: none">• disabled - the interface will not use ARP• enabled - the interface will use ARP• proxy-arp - the interface will use the ARP proxy feature• reply-only - the interface will only reply to requests originated from matching IP address/MAC address combinations which are entered as static entries in the "/ip arp" table. No dynamic entries will be automatically stored in the "/ip arp" table. Therefore for communications to be successful, a valid static entry must already exist.
arp-interval (<i>time</i> ; Default: 00:00:00.100)	Time in milliseconds defines how often to monitor ARP requests
arp-ip-targets (<i>IP address</i> ; Default:)	IP target address which will be monitored if link-monitoring is set to arp. You can specify multiple IP addresses, separated by a comma
comment (<i>string</i> ; Default:)	Short description of the interface
disabled (<i>yes no</i> ; Default: no)	Changes whether the bonding interface is disabled
down-delay (<i>time</i> ; Default: 00:00:00)	If a link failure has been detected, the bonding interface is disabled for a down-delay time. The value should be a multiple of mii-interval, otherwise, it will be rounded down to the nearest value. This property only has an effect when <code>link-monitoring</code> is set to <code>mii</code> .
forced-mac-address (<i>MAC address</i> ; Default: none)	By default, the bonding interface will use the MAC address of the first selected slave interface. This property allows to configure static MAC address for the bond interface (all zeros, broadcast or multicast addresses will not apply). RouterOS will automatically change the MAC address for slave interfaces and it will be visible in <code>/interface ethernet</code> configuration export
lACP-rate (<i>1sec 30secs</i> ; Default: 30secs)	Link Aggregation Control Protocol rate specifies how often to exchange with LACPDUs between bonding peers. Used to determine whether a link is up or other changes have occurred in the network. LACP tries to adapt to these changes providing failover.
lACP-user-key (<i>integer: 0..1023</i> ; Default: 0)	Specifies the upper 10 bits of the port key. The lower 6 bits are automatically assigned based on individual port link speed and duplex. The setting is available only since RouterOS v7.3.

Property	Description
link-monitoring (<i>arp mii none</i> ; Default: mii)	<p>Method to use for monitoring the link (whether it is up or down)</p> <ul style="list-style-type: none"> • arp - uses Address Resolution Protocol to determine whether the remote interface is reachable • mii - uses Media Independent Interface to determine link status. Link status determination relies on the device driver. • none - no method for link monitoring is used. <p>Note: some bonding modes require specific link monitoring to work properly.</p>
min-links (<i>integer: 0..4294967295</i> ; Default: 0)	<p>How many active slave links needed for bonding to become active</p>
mii-interval (<i>time</i> ; Default: 00:00:00.100)	<p>How often to monitor the link for failures (the parameter used only if link-monitoring is mii)</p>
mlag-id (<i>integer: 0..4294967295</i> ; Default:)	<p>Changes MLAG ID for bonding interface. The same MLAG ID should be used on both peer devices to successfully create a single MLAG. See more details on MLAG.</p>

Property	Description
<p>mode (<i>802.3ad</i> <i>active-backup</i> <i>balance-alb</i> <i>balance-rr</i> <i>balance-tlb</i> <i>balance-xor</i> <i>broadcast</i>; Default: balance-rr)</p>	<p>Specifies one of the bonding policies</p> <ul style="list-style-type: none"> • 802.3ad - IEEE 802.3ad dynamic link aggregation. In this mode, the interfaces are aggregated in a group where each slave shares the same speed. It provides fault tolerance and load balancing. Slave selection for outgoing traffic is done according to the transmit-hash-policy more> • active-backup - provides link backup. Only one slave can be active at a time. Another slave only becomes active, if the first one fails. more> • balance-alb - adaptive load balancing. The same as balance-tlb but received traffic is also balanced. The device driver should have support for changing it's MAC address. more> • balance-rr - round-robin load balancing. Slaves in a bonding interface will transmit and receive data in sequential order. It provides load balancing and fault tolerance. more> • balance-tlb - Outgoing traffic is distributed according to the current load on each slave. Incoming traffic is not balanced and is received by the current slave. If receiving slave fails, then another slave takes the MAC address of the failed slave. more> • balance-xor - Transmit based on the selected transmit-hash-policy. This mode provides load balancing and fault tolerance. more> • broadcast - Broadcasts the same data on all interfaces at once. This provides fault tolerance but slows down traffic throughput on some slow machines. more>
<p>mtu (<i>integer</i>; Default: 1500)</p>	<p>Maximum Transmit Unit in bytes. Must be smaller or equal to the smallest L2MTU value of a bonding slave. L2MTU of a bonding interface is determined by the lowest L2MTU value among its slave interfaces</p>
<p>name (<i>string</i>; Default:)</p>	<p>Name of the bonding interface</p>
<p>primary (<i>string</i>; Default: none)</p>	<p>Controls the primary interface between active slave ports, works only for active-backup, balance-tlb and balance-alb modes. For active-backup mode, it controls which running interface is supposed to send and receive the traffic. For balance-tlb mode, it controls which running interface is supposed to receive all the traffic, but for balance-alb mode, it controls which interface is supposed to receive the unbalanced traffic (the non-IPv4 traffic). When none of the interfaces are selected as primary, device will automatically select the interface that is configured as the first one.</p>

Property	Description
slaves (<i>string</i> ; Default: none)	At least two ethernet-like interfaces separated by a comma, which will be used for bonding
up-delay (<i>time</i> ; Default: 00:00:00)	If a link has been brought up, the bonding interface is disabled for up-delay time and after this time it is enabled. The value should be a multiple of mii-interval, otherwise, it will be rounded down to the nearest value. This property only has an effect when <code>link-monitoring</code> is set to <code>mii</code> .
transmit-hash-policy (<i>layer-2 layer-2-and-3 layer-3-and-4</i> ; Default: layer-2)	<p>Selects the transmit hash policy to use for slave selection in balance-xor and 802.3ad modes</p> <ul style="list-style-type: none"> • layer-2 - Uses XOR of hardware MAC addresses to generate the hash. This algorithm will place all traffic to a particular network peer on the same slave. This algorithm is 802.3ad compliant. • layer-2-and-3 - This policy uses a combination of layer2 and layer3 protocol information to generate the hash. Uses XOR of hardware MAC addresses and IP addresses to generate the hash. This algorithm will place all traffic to a particular network peer on the same slave. For non-IP traffic, the formula is the same as for the layer2 transmit hash policy. This policy is intended to provide a more balanced distribution of traffic than layer2 alone, especially in environments where a layer3 gateway device is required to reach most destinations. This algorithm is 802.3ad compliant. • layer-3-and-4 - This policy uses upper layer protocol information, when available, to generate the hash. This allows for traffic to a particular network peer to span multiple slaves, although a single connection will not span multiple slaves. For fragmented TCP or UDP packets and all other IP protocol traffic, the source and destination port information is omitted. For non-IP traffic, the formula is the same as for the layer2 transmit hash policy. This algorithm is not fully 802.3ad compliant.

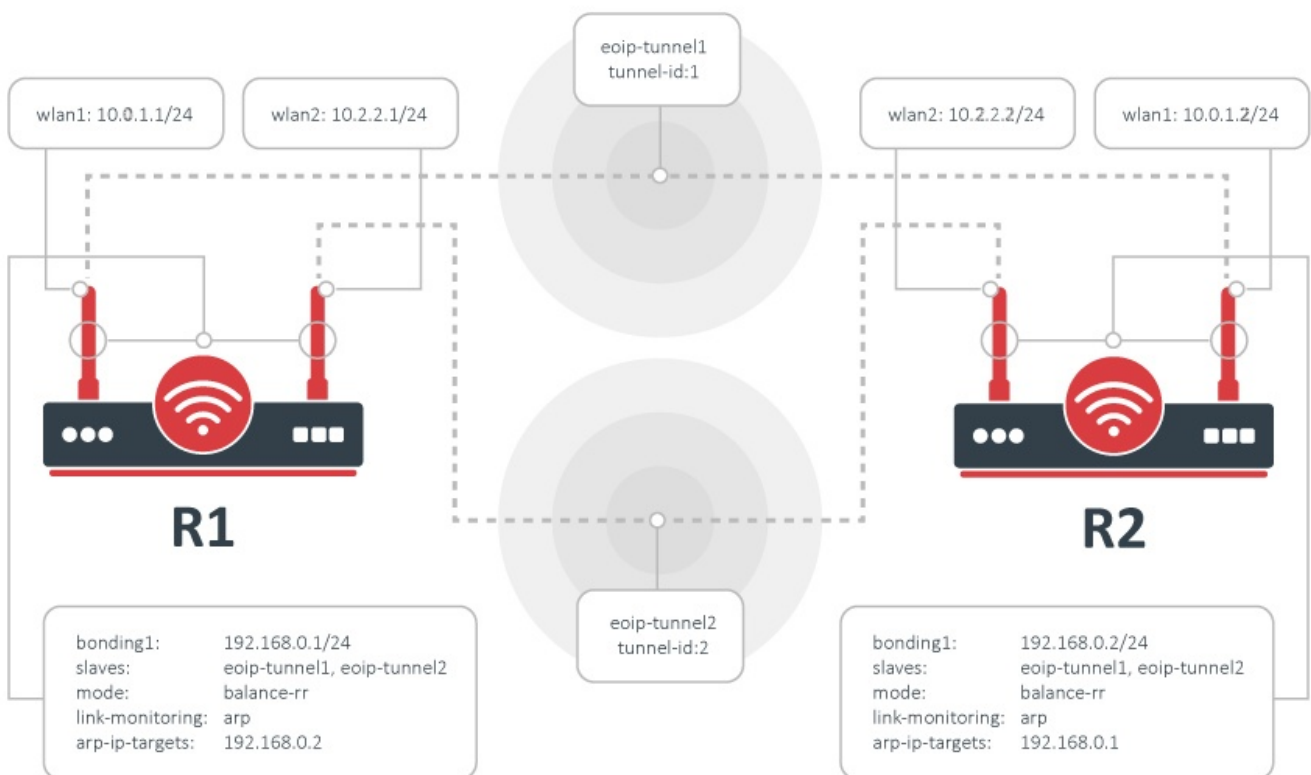
Bonding Examples

Bonding EoIP tunnels over two wireless links

This is an example of aggregating multiple network interfaces into a single pipe. In particular, it is shown how to aggregate multiple virtual (EoIP) interfaces to get maximum throughput (MT) with emphasis on availability.

Network Diagram

Two routers R1 and R2 are interconnected via wireless links. Wireless interfaces on both sides have assigned IP addresses.



Configuration

Bonding could be used only on OSI layer 2 (Ethernet level) connections. Thus we need to create EoIP interfaces on each of the wireless links. This is done as follows:

on router R1:

```
/interface eoip add remote-address=10.0.1.1/24 tunnel-id=1
/interface eoip add remote-address=10.2.2.1/24 tunnel-id=2
```

and on router R2:

```
/interface eoip add remote-address=10.0.1.2/24 tunnel-id=1
/interface eoip add remote-address=10.2.2.2/24 tunnel-id=2
```

The second step is to add a bonding interface and specify EoIP interfaces as slaves:

R1:

```
/interface bonding add slaves=eoip-tunnel1,eoip-tunnel2 mode=balance-rr
```

R2:

```
/interface bonding add slaves=eoip-tunnel1,eoip-tunnel2 mode=balance-rr
```

The last step is to add IP addresses to the bonding interfaces:

R1:

```
/ip address add address 192.168.0.1/24 interface=bonding1
```

R2:

```
/ip address add address 192.168.0.2/24 interface=bonding1
```

Test the configuration

Now two routers are able to reach each other using addresses from the 192.168.0.0/24 network. To verify bonding interface functionality, do the following:

R1:

```
/interface monitor-traffic eoip-tunnel1,eoip-tunnel2
```

R2:

```
/tool bandwidth-test 192.168.0.1 direction=transmit
```

You should see that traffic is distributed equally across both EoIP interfaces:

```
/int monitor-traffic eoip-tunnel1,eoip-tunnel2
received-packets-per-second: 685    685
received-bits-per-second: 8.0Mbps  8.0Mbps
sent-packets-per-second: 21     20
sent-bits-per-second: 11.9kbps 11.0kbps
```

```
received-packets-per-second: 898 899
received-bits-per-second: 10.6Mbps 10.6Mbps
sent-packets-per-second: 20 21
sent-bits-per-second: 11.0kbps 11.9kbps
received-packets-per-second: 975 975
received-bits-per-second: 11.5Mbps 11.5Mbps
sent-packets-per-second: 22 22
sent-bits-per-second: 12.4kbps 12.3kbps
received-packets-per-second: 980 980
received-bits-per-second: 11.6Mbps 11.6Mbps
sent-packets-per-second: 21 21
sent-bits-per-second: 11.9kbps 11.8kbps
received-packets-per-second: 977 977
received-bits-per-second: 11.6Mbps 11.5Mbps
sent-packets-per-second: 21 21
sent-bits-per-second: 11.9kbps 11.8kbps
-- [Q quit|D dump|C-z pause]
```

Link Monitoring

It is easy to notice that with the configuration above as soon as any individual link fails, the bonding interface throughput collapses. That's because no link monitoring is performed, consequently, the bonding driver is unaware of problems with the underlying links. Enabling link monitoring is a must in most bonding configurations. To enable ARP link monitoring, do the following:

R1:

```
/interface bonding set bonding1 link-monitoring=arp arp-ip-targets=192.168.0.2
```

R2:

```
/interface bonding set bonding1 link-monitoring=arp arp-ip-targets=192.168.0.1
```